Securing Data Transfers:

An Integrity Algorithm for Error Recovery Triangulation

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# Abstract

*Transferring data is one of the key operations performed by millions of users every day. Users do this by issuing direct commands, such as file transfer commands, or indirectly as a feature invoked by numerous end-user applications. The most important security characteristic of a successful data exchange is the integrity of that data. The receiver user desires to acquire data that has not been modified through malicious acts, or simple human or machine error. Applications that rely on the Transfer Control Protocol (TCP) as the main mechanism to provide end-to-end reliability, including error and sequence control, do not check the integrity of the file being transmitted prior to the transfer. In this paper, we present an overview of current data transfer mechanisms and their security provisions and propose an internal integrity mechanism that provides a triangulation means of error control through the use of one-way hash functions based on the original file being transferred; and a discussion of the implications and limitations that such a mechanism imparts on data transfer mechanisms.*

**Keywords**: Secure Data Transfers, Security, Integrity, Error Recovery, Triangulation, Hash Product.

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# 1. Introduction

Data transfers constitute one of the most commonly performed tasks in the Internet today. Users publishing files to a website, moving files that are too large for an email application, transferring files securely between different computers, uploading photos to services such as Flickr (www.flickr.com) or backing up a to a remote server are all engaged in data transfer operations. File Transfer Protocol (FTP), for example, is still used to perform bulk data transfers across networks (Grzywa et al., 2001). It is important to note that this simple-sounding task (data transfer) is essential for the interoperation of networked computers, especially when the different computing environments of modern heterogeneous networks are taken into account.

Potential user devices in the network range from Personal Digital Assistants (PDAs) and mobile phones to specialized servers and large-scale corporate systems. All of these devices could operate using different word formats; they might store the data in different forms and forward packets in non-compatible ways (big-endian vs. little-endian computers) (Tanenbaum, 2003).

# 2. Overview of Data Transfer Mechanisms

There are two main approaches to perform data transfers in today’s IP-based networks. One is based on FTP and several variations to that protocol and the second, and much newer, approach is based on peer-to-peer networking configurations. Let’s review both of them.

## 2.1 FTP Overview

Essentially, FTP is a client-server protocol that facilitates the transfer of files between two computers connected via a network, such as the Internet (Brown & Jaatun, 1992). The protocol has several fundamental components that perform the required functions for establishing the sessions between systems, and coordinates the transfer of the files. These components are as follows:

* Server-FTP that consists of the Server Protocol Interpreter and the Server Data Transfer Process;
* Server-side File System;
* User-FTP that consists of:
  + the User Interface,
  + User Protocol Interpreter, and
  + the User Data Transfer Process;
* User-side File System;
* A User; and
* A communication channel for the control and data connections.

An alternative configuration of FTP provides for the transfer between two servers through a third system that provides the control function between the servers. In both of the configurations, FTP requires that the control channel be open during the data transfer process (Postel & Reynolds, 1985).

The FTP architecture is displayed below (see Figure 1):

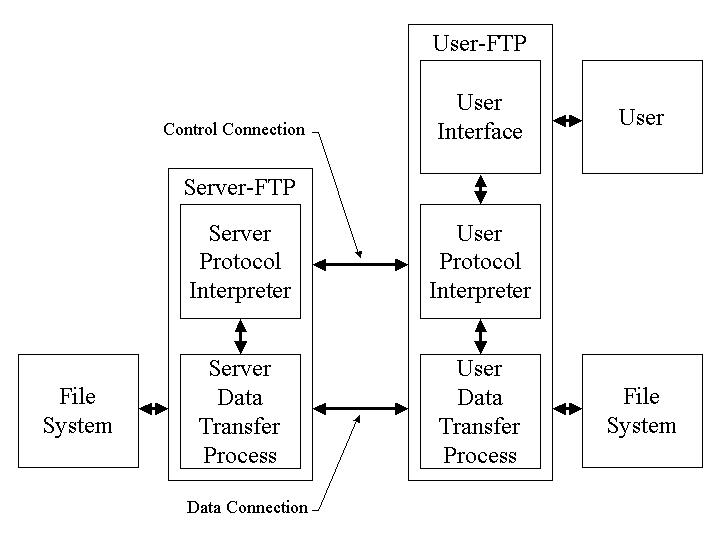


Figure 1: FTP Architecture

(Source: Postel & Reynolds, 1985)

### 2.1.1 Telnet Authentication and Encryption

As previously discussed, FTP utilizes the control connection to coordinate the data connection and execute commands on the File Systems of both the user and the server. FTP utilizes the Telnet protocol to execute the commands (Postel & Reynolds, 1985). This fundamental design lends itself to security breaches that may permit eavesdropping of user ID’s, passwords, file names, and other information passed through the control channel. It also may allow an active attacker to change settings and execute commands on the file system (Brown & Jaatun, 1992). This fundamental security flaw was initially addressed when Borman (1993) proposed the passing of authentication information, and a mechanism to enable encryption of the data after successful authentication for the Telnet protocol. The result was that user passwords would not be sent in clear text, and the data stream would be encrypted utilizing any general authentication and encryption system. However, it should noted that the “Telnet authentication and encryption option does not provide for integrity protection only (without confidentiality), and does not address the protection of the data channel” (Horowitz & Lunt, 1997).

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| Scenarios | | | | |
| Send Request Response | Request Hash | Storage Hash | Generated Hash | Action Required / Notification Returned |
| h1 | h1 | h1 | Validated file, integrity confirmed, send file |
| h1 | h1 | h2 | File has changed since original message, original hash may not have been updated, confirm file integrity before continuing, regenerate hash and reconfirm |
| h1 | h2 | h2 | Message to Recipient is invalid or out of date, may be invalid request, request retransmission of message |
| h1 | h2 | h1 | Original hash may be corrupted, update hash file, confirm before transmitting file |
| h2 | h1 | h1 | Message to Recipient is invalid or out of date, issue new notification, request retransmission of message |
| h2 | h2 | h1 | File has changed since original message, original hash may not have been updated, confirm file integrity before continuing, regenerate hash and reconfirm |
| h2 | h1 | h2 | Original hash may be corrupted, update hash, confirm before transmitting file |
| h2 | h1 | h3 | Re-evaluate all controls for file, discontinue distribution of file, and  everyone is going to hell in a hand-basket ☺ |

**Table 1**: Send Request Response

# 5. Conclusion

A basic requirement by users is that files be transferred efficiently and without modification. File integrity, therefore, becomes a critical element in the sharing of files. During the transfer of a file, a self-correcting system would permit the user to acquire a file while the error controls would operate transparently in the background. An added feature of the Hash Triplet is that it provides the system with specific points of origins for the root error causes and allows system designers to promote a wider range of remedies.

## 5.1 Formatting References

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## Appendix 1

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